Epigraphs, Infimal convolution & Moreau-Yosida envelope

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Epigraphs

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epi
$$f \coloneqq \left\{ (\boldsymbol{x}, \alpha) \in \mathbb{R}^n \times \mathbb{R} : \alpha \ge f(\boldsymbol{x}) \right\}.$$

Graph and epigraph

► (Graph of abstract function) Given two sets X (domain) and Y (codomain), the graph of the function f : X → Y is the set

$$\operatorname{graph} f \coloneqq \Big\{ (\boldsymbol{x}, \boldsymbol{y}) \in X \times Y : \boldsymbol{y} = f(\boldsymbol{x}) \Big\}.$$

• (Epigraph of function from \mathbb{R}^n to $\overline{\mathbb{R}}$) Let $\overline{\mathbb{R}} := \mathbb{R} \cup \{+\infty\}$ be the extended reals. Now consider $X = \mathbb{R}^n$ and $Y = \overline{\mathbb{R}}$. For $f : \mathbb{R}^n \to \overline{\mathbb{R}}$, the epigraph of f is the set

$$\operatorname{epi} f \coloneqq \Big\{ (\boldsymbol{x}, \alpha) \in \mathbb{R}^n \times \mathbb{R} : \alpha \ge f(\boldsymbol{x}) \Big\}.$$

▶ (Strict epigraph) $epi_S f = epi f \setminus graph f$, or equivalently

$$\operatorname{epi}_{S} f \coloneqq \Big\{ (\boldsymbol{x}, \alpha) \in \mathbb{R}^{n} \times \mathbb{R} : \alpha > f(\boldsymbol{x}) \Big\}.$$

Remarks

$$\operatorname{epi} f \coloneqq \Big\{ ({\boldsymbol x}, \alpha) \in \mathbb{R}^n \times \mathbb{R} \, : \, \alpha \geq f({\boldsymbol x}) \Big\}$$

• $f: \mathbb{R}^n \to \overline{\mathbb{R}}$ but $\operatorname{epi} f$ is defined to be a subset of $\mathbb{R}^n \times \mathbb{R}$, not $\mathbb{R}^n \times \overline{\mathbb{R}}$.

• This is intentional to define epi f as a subset of $\mathbb{R}^n \times \mathbb{R}$.

- \mathbb{R}^n is a vector space
- $\mathbb{R}^n \times \mathbb{R}$ is a vector space
- $\mathbb{R}^n imes \overline{\mathbb{R}}$ is not a vector space: \nexists additive identity for $\infty + \infty$

(Being a vector space allows to use tools from real analysis and functional analysis.)

• (At infinity) If $f(x_0) = +\infty$ at $x = x_0$, then $(x_0, +\infty) \notin epi f$.

Two extreme cases

• If
$$f_{\infty^+}(\boldsymbol{x}) = +\infty \ \forall \boldsymbol{x}$$
, then $\operatorname{epi} f_{\infty^+}$ is \varnothing .

• If
$$f_{\infty^-}(x) = -\infty \forall x$$
, then $\operatorname{epi} f_{\infty^-}$ is the whole $\mathbb{R}^n \times \mathbb{R}$
(Convention: $-\infty < +\infty$)

Empty epigraph Whole space epigraph Visualization of graph f and epi f

- $\operatorname{epi} f = \operatorname{all}$ the points of \mathbb{R}^{n+1} lying on or above $\operatorname{graph} f$.
- ► Example: $f(x) = x^2$
 - ▶ n = 1 (1-dimensional)
 - graph $f \coloneqq \{(x, y) \in \mathbb{R} \times \mathbb{R} : y = f(x)\}$ is a 1d curve in a 2d space.
 - $\operatorname{epi} f := \left\{ (x, \alpha) \in \mathbb{R} \times \mathbb{R} : \alpha \ge f(x) \right\}$ is a 2d set in a 2d space.



Level sets: a concept related to epigraph.

$$\begin{split} \operatorname{lev}_{\leq \alpha} f &\coloneqq \left\{ \boldsymbol{x} \in \mathbb{R}^n \,:\, f(\boldsymbol{x}) \leq \alpha \right\} \\ \operatorname{lev}_{<\alpha} f &\coloneqq \left\{ \boldsymbol{x} \in \mathbb{R}^n \,:\, f(\boldsymbol{x}) < \alpha \right\} \\ \blacktriangleright & \operatorname{For} f : \mathbb{R}^n \to \mathbb{R} \text{ and } \alpha \in \mathbb{R}, \text{ define } \begin{split} \operatorname{lev}_{=\alpha} f &\coloneqq \left\{ \boldsymbol{x} \in \mathbb{R}^n \,:\, f(\boldsymbol{x}) < \alpha \right\} \\ \operatorname{lev}_{>\alpha} f &\coloneqq \left\{ \boldsymbol{x} \in \mathbb{R}^n \,:\, f(\boldsymbol{x}) > \alpha \right\} \\ \operatorname{lev}_{\geq \alpha} f &\coloneqq \left\{ \boldsymbol{x} \in \mathbb{R}^n \,:\, f(\boldsymbol{x}) > \alpha \right\} \\ \operatorname{lev}_{\geq \alpha} f &\coloneqq \left\{ \boldsymbol{x} \in \mathbb{R}^n \,:\, f(\boldsymbol{x}) \geq \alpha \right\} \end{split}$$

The important one for minimization is $lev_{\leq \alpha} f$ (named sublevel sets).

- ► lev is a subset of domain, not codomain.
- If $\alpha = \inf f$, then $\operatorname{lev}_{\leq \alpha} f = \operatorname{lev}_{=\alpha} f = \operatorname{argmin} f$.
- Level sets can be
 - empty: $lev_{\leq -1}(x^2)$, no x makes $x^2 \leq -1$
 - non-continuous: $lev_{=0} \sin(x)$, then $x = \{n\pi\}_{n \in \mathbb{N}}$ is a set of dots
 - ▶ non-convex: $lev_{=0} sin(x)$

An illustration of $lev_{\leq \alpha} f$



Picture from Rockafellar, R. Tyrrell, and Roger J-B. Wets. Variational analysis. Springer, 2009.

Optimization newbie: "Why talk about epigraph?"

Optimization expert: "it is useful!"

The main idea

- Many properties of f has a counterpart in epi f.
- Sometimes it is easier to work with epi f than with f.

Properties of f

Properties of $\operatorname{epi} f$

epi f is nonempty

epi f is a closed set epi f is nonempty and closed in $\mathbb{R}^n \times \mathbb{R}$ epi f is a convex set epi $_S f$ is a convex set

Minkowski sum of epi f and epi g

► *f* is proper

- f is closed
- f is lower semicontinuous on \mathbb{R}^n
- f is convex
- f is strictly convex
- $f \Box g$ (Infimal convolution of f and g)

Example: f is convex $\iff epi f$ is convex

(To avoid confusing the bracket of graph and bracket of a number, here we use $\{x_1, y_1\} \in {
m epi}\, f)$

► (⇐=)

▶ epi f is a convex set $\iff \{x_1, y_1\} \in epi f$, $\{x_2, y_2\} \in epi f$ implies for $t \in [0, 1]$ we have

$$(1-t)\left\{\boldsymbol{x}_{1}, y_{1}\right\} + t\left\{\boldsymbol{x}_{2}, y_{2}\right\} = \left\{(1-t)\boldsymbol{x}_{1} + t\boldsymbol{x}_{2}, (1-t)y_{1} + ty_{2}\right\} \in \operatorname{epi} f.$$
 (***)

By definition of epigraph, (* * *) is equivalent to

$$f(\boldsymbol{x}_1) \leq y_1 \text{ and } f(\boldsymbol{x}_2) \leq y_2 \quad \text{implies} \quad f\Big((1-t)\boldsymbol{x}_1 + t\boldsymbol{x}_2\Big) \leq (1-t)y_1 + ty_2$$

Choose $f(\boldsymbol{x}_1) \leq y_1$ and $f(\boldsymbol{x}_2) \leq y_2$ gives $f((1-t)\boldsymbol{x}_1 + t\boldsymbol{x}_2) \leq (1-t)f(\boldsymbol{x}_1) + tf(\boldsymbol{x}_2)$, so f is convex.

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$$(f_1 \Box f_2)(\boldsymbol{x}) \coloneqq \inf_{\boldsymbol{x}_1 + \boldsymbol{x}_2 = \boldsymbol{x}} f_1(\boldsymbol{x}_1) + f_2(\boldsymbol{x}_2)$$

Infimal convolution

• Given $f_1 : \mathbb{R}^n \to \overline{\mathbb{R}}$ and $f_2 : \mathbb{R}^n \to \overline{\mathbb{R}}$. The infimal convolution $f_1 \Box f_2 : \mathbb{R}^n \to \overline{\mathbb{R}}$ is

$$egin{array}{rll} (f_1 \Box f_2)(m{x}) &\coloneqq& \inf_{m{x}_1 + m{x}_2 = m{x}} f_1(m{x}_1) + f_2(m{x}_2) \ &=& \inf_{m{x}_1} & f_1(m{x}_1) + f_2(m{x} - m{x}_1) \ &=& \inf_{m{x}_2} & f_1(m{x} - m{x}_2) + f_2(m{x}_2) \end{array}$$

Convention: $\infty - \infty = \infty$ and $\inf \varnothing = +\infty$

History

Earliest(?) work

Hausdorff, Uber halbstetige Funktionen und deren Verallgemeinerung, Math. Zeit. 5 (1919), 292-309.

Fenchel, "Convex Cones, Sets, and Functions", Lecture Notes, Princeton University, Princeton, 1953.

First systematic study of infimal convolution Moreau, Inf-convolution, Sém. d'Math. Montpellier (1963), 3.1-3.48

Later works by Attouch, Rockafellar, Hiriart-Urruty, etc

Thomas Stromberg's PhD thesis (1994): a nice summary.

How infimal convolution gets its name

$$(f_1 \Box f_2)(\boldsymbol{x}) := \inf_{\boldsymbol{x}_1} f_1(\boldsymbol{x}_1) + f_2(\boldsymbol{x} - \boldsymbol{x}_2).$$

Definition of (integral) convolution

(Examples: Laplace transform, Fourier transform.)

$$(f*g)(t) := \int_{\infty}^{\infty} f(\tau)g(t-\tau)d\tau.$$

•
$$f_1(\boldsymbol{x}_1) + f_2(\boldsymbol{x} - \boldsymbol{x}_2)$$
 "looks similar" to (integral) convolution

- There is $\inf_{w} \implies$ people name it infimal.
- Deep fact: integral convolution is in $(+, \times)$ -algebra
 - ► Integration = summation
 - You combine f and g by multiplication

Infimal convolution is in $(\min, +)$ -algebra (tropical semi-ring)

- ► The summation is replaced by min
- You combine f and g by addition

What infimal convolution solves: an economics example

- ▶ You want to buy totally n hamburgers, from MacDonald and Burger King. Suppose buying n_1 hamburgers from MacDonald costs you $f(n_1)$, and if you buy n_2 hamburgers from Burger King, the price is $g(n_2)$.
- Vou want to find the infimum of the total cost $f(n_1) + g(n_2)$ subject to the constraint $n_1 + n_2 = n$. I.e., you want to find the "cheapest way" to buy n hamburgers.
- This problem is exactly: calculate $(f \Box g)(n)$

$$f\Box g)(n) = \inf_{n_1+n_2=n} f(n_1) + g(n_2) = \underbrace{\inf_{n_1} f(n_1) + g(n-n_1)}_{\text{focus on } n_1} = \underbrace{\inf_{n_2} f(n-n_2) + g(n_2)}_{\text{focus on } n_2}.$$

this also means infimal convolution is commutative

Infimal convolution is commutative: $f \Box g = g \Box f$.

Example. Infimal convolution of two indicator functions

$$\begin{array}{rcl} f_1 \Box f_2)({\bm x}) &\coloneqq & \inf_{{\bm x}_1 + {\bm x}_2 = {\bm x}} f_1({\bm x}_1) + f_2({\bm x}_2) \\ &= & \inf_{{\bm x}_1} & f_1({\bm x}_1) + f_2({\bm x} - {\bm x}_1) \\ &= & \inf_{{\bm x}_2} & f_1({\bm x} - {\bm x}_2) + f_2({\bm x}_2) \end{array}$$

• Given two sets C_1, C_2 and two indicator functions i_{C_1}, i_{C_2} .

$$(i_{C_1} \Box i_{C_2})(\boldsymbol{x}) = \inf_{\boldsymbol{x}_1 + \boldsymbol{x}_2 = \boldsymbol{x}} i_{C_1}(\boldsymbol{x}_1) + i_{C_2}(\boldsymbol{x}_2) = i_{C_1 \oplus C_2}(\boldsymbol{x})$$

 \oplus is Minkowski sum of sets: $P \oplus Q \coloneqq \{p+q \mid p \in P, q \in Q\}.$

- Minkowski sum keeps convexity of sets, so $C_1 \oplus C_2$ is a convex set and $i_{C_1 \oplus C_2}$ is a convex function. Here we see that inf-convolution of two convex functions is a convex function.
- ▶ In general, if $f_1 : \mathbb{R}^n \to \overline{\mathbb{R}}$ and $f_2 : \mathbb{R}^n \to \overline{\mathbb{R}}$ are convex, then $f_1 \Box f_2$ is also convex. Proof. We prove inf-convolution preserves convexity using definition and operations that preserve convexity. By definition, $(f_1 \Box f_2)(x) = \inf_{x_1} h(x, x_1)$ where $h(x, x_1) = f_1(x_1) + f_2(x - x_1)$. By assumption $f_1(x_1)$ is convex and $f_2(x_1)$ is convex. The function $f_2(x - x_1)$ is $f_2(x_1)$ with argument x_1 under a translation to $x - x_1$ so $f_2(x - x_1)$ is convex. Now $h(x, x_1)$ is the sum of two convex functions on x_1 , thus it is convex.

Infimal convolution is also called epi-addition

$$(f_1 \Box f_2)(\boldsymbol{x}) := \inf_{\boldsymbol{x}_2} f_1(\boldsymbol{x} - \boldsymbol{x}_2) + f_2(\boldsymbol{x}_2).$$

- $\operatorname{epi}_{S}(f_{1} \Box f_{2}) = \operatorname{epi}_{S} f_{1} \oplus \operatorname{epi}_{S} f_{2} \iff \operatorname{inf-convolution} of \operatorname{convex} functions is \operatorname{convex}^{1}$
- $\operatorname{epi}(f_1 \Box f_2) \supseteq \operatorname{epi} f_1 \oplus \operatorname{epi} f_2$
- ▶ epi $(f_1 \Box f_2)$ = epi $f_1 \oplus$ epi f_2 if inf-convolution is exact Exact means the inf is gone: $(f_1 \Box f_2)(x) = f_1(x - x_2) + f_2(x_2)$.
- For proof, see Jean Jacques Moreau. Inf-convolution, sous-additivité, convexité des fonctions numériques. Journal de Mathématiques Pures et Appliquées, 1970. https://hal.archives-ouvertes.fr/hal-02162006

¹Remark 2.3.3 in Urruty, Jean-Baptiste Hiriart, and Claude Lemaréchal. Convex analysis and minimization algorithms. Springer-Verlag, 1993

The proof of $epi(f_1 \Box f_2) \supseteq epi f_1 \oplus epi f_2$

- ► Take $\{x, \alpha\} \in \text{epi } f_1 \oplus \text{epi } f_2$. Since the element $\{x, \alpha\}$ is inside the sum of two sets (epi f_1 and epi f_2), that means we can decompose $\{x, \alpha\}$ as the sum of element from each set. I.e., we have $x = x_1 + x_2$ and $\alpha = \alpha_1 + \alpha_2$ that $\{x_1, \alpha_1\} \in \text{epi } f_1$ and $\{x_2, \alpha_2\} \in \text{epi } f_2$. (It means given $\{x, \alpha\}$, there is exist the pair $\{x_1, \alpha_1\}, \{x_2, \alpha_2\}$ that fulfill the above conditions)
- implies $f_1(\boldsymbol{x}_1) \leq \alpha_1$, implies $f_2(\boldsymbol{x}_2) \leq \alpha_2$ and $f_1(\boldsymbol{x}_1) + f_2(\boldsymbol{x}_2) \leq \alpha_1 + \alpha_2 = \alpha$.
- Now consider $(f_1 \Box f_2)(x) \coloneqq \inf_{\substack{y_1 + y_2 = x}} f_1(y_1) + f_2(y_2)$. As $x_1 + x_2 = x = y_1 + y_2$, the infimum $\inf_{\substack{y_1 + y_2 = x}}$ is the smallest among all pair that sum to x, so

$$(f_1 \Box f_2)(\boldsymbol{x}) \coloneqq \inf_{\boldsymbol{y}_1 + \boldsymbol{y}_2 = \boldsymbol{x}} f_1(\boldsymbol{y}_1) + f_2(\boldsymbol{y}_2) \le f_1(\boldsymbol{x}_1) + f_2(\boldsymbol{x}_2) \le \alpha.$$

So $\{x, \alpha\} \in \operatorname{epi}(f_1 \Box f_2)$.

• What we just showed is $\{x, \alpha\} \in epi f_1 \oplus epi f_2 \implies \{x, \alpha\} \in epi (f_1 \Box f_2)$, so in set language we have $epi (f_1 \Box f_2) \supseteq epi f_1 \oplus epi f_2$.

Example. Pictorial illustration of epi-addition (and Minkowski sum)

"Epi-addition: sliding the blue curve on red curve and perform union operation gives the black curve" $% \left({{{\mathbf{r}}_{i}}} \right) = {{\mathbf{r}}_{i}} \right)$

What about $f_1 \Box f_2 \Box f_3$? Inf-convolution is associative

$$\begin{split} f_1 \Box (f_2 \Box f_3)(t) &= \inf_{x+y=t} \left\{ f_1(x) + (f_2 \Box f_3)(y) \right\} & \text{by definition} \\ &= \inf_{x+y=t} \left\{ f_1(x) + \left\{ \inf_{z+w=y} f_2(w) + f_3(z) \right\} \right\} & \text{by definition} \\ &= \inf_{x+y=t} \left\{ f_1(x) + f_2(w) + f_3(z) \right\} & \text{you can move inf} \\ &= \inf_{x+z+w=t} \left\{ f_1(x) + f_2(w) + f_3(z) \right\} & \text{combine } x + y = t, z + w = y \\ &= \inf_{\substack{r+z=t \\ x+w=r}} \left\{ f_1(x) + f_2(w) + f_3(z) \right\} & \text{let } t = r + z, x + w = r \\ &= \inf_{\substack{r+z=t \\ r+z=t}} \left\{ \left\{ \inf_{x+w=r} f_1(x) + f_2(w) \right\} + f_3(z) \right\} & \text{you can move inf} \\ &= \inf_{\substack{r+z=t \\ r+z=t}} \left\{ (f_1 \Box f_2)(r) + f_3(z) \right\} & \text{by definition} \\ &= (f_1 \Box f_2) \Box f_3(t) & \text{by definition} \end{split}$$

Properties of inf-convolution

We already see

- $\blacktriangleright \ f \Box g = g \Box f$
- $\blacktriangleright \ f \Box g \Box h = (f \Box g) \Box h = f \Box (g \Box h)$
- $\blacktriangleright \ f,g \text{ convex } \implies f \Box g \text{ convex}$

Useful table

f	g	$f\Box g$
f	0	$\inf_{x} f(x)$
i_C	$\ \cdot\ _2$	d_C
i_{C_1}	i_{C_2}	$i_{C_1\oplus C_2}$
f	i_x	$f(\cdot - x)$
f	$\langle s,\cdot angle$	$\langle s, \cdot \rangle - f^*(s)$
f convex	f	$2f\left(\frac{\cdot}{2}\right)$

commutative associative

inf-convolution preserves convexity

Distance function
$$d_C(x) = \inf_{c \in C} ||x - c||_2$$

► Indicator function
$$i_C(x) = \begin{cases} 0 & x \in C \\ +\infty & x \notin C \end{cases}$$

• Conjugate
$$f^*(x) = \sup_u \langle u, x \rangle - f(x)$$

Inf-convolution and conjugate: $(f \Box g)^* = f^* + g^*$ Let $f, g : \mathbb{R}^n \to \overline{\mathbb{R}}$ be proper functions.

$$\begin{aligned} (f\Box g)^*(\boldsymbol{y}) &= \sup_{\boldsymbol{x}} \left\{ \langle \boldsymbol{y}, \boldsymbol{x} \rangle - (f\Box g)(\boldsymbol{x}) \right\} & \text{by definition of conjugate} \\ &= \sup_{\boldsymbol{x}} \left\{ \langle \boldsymbol{y}, \boldsymbol{x} \rangle - \inf_{\boldsymbol{u}} \left[f(\boldsymbol{u}) + g(\boldsymbol{x} - \boldsymbol{u}) \right] \right\} & \text{by definition of inf-convolution} \\ &= \sup_{\boldsymbol{x}} \sup_{\boldsymbol{u}} \left\{ \langle \boldsymbol{y}, \boldsymbol{x} \rangle - \left[f(\boldsymbol{u}) + g(\boldsymbol{x} - \boldsymbol{u}) \right] \right\} & -\inf_{\boldsymbol{u}} = +\sup_{\boldsymbol{u}} \\ &= \sup_{\boldsymbol{x}, \boldsymbol{u}} \left\{ \langle \boldsymbol{y}, \boldsymbol{x} - \boldsymbol{u} + \boldsymbol{u} \rangle - f(\boldsymbol{u}) - g(\boldsymbol{x} - \boldsymbol{u}) \right\} \\ &= \sup_{\boldsymbol{x}, \boldsymbol{u}} \left\{ \langle \boldsymbol{y}, \boldsymbol{x} - \boldsymbol{u} + \boldsymbol{u} \rangle - f(\boldsymbol{u}) - g(\boldsymbol{x} - \boldsymbol{u}) \right\} \\ &= \sup_{\boldsymbol{x}, \boldsymbol{u}} \left\{ \langle \boldsymbol{y}, \boldsymbol{x} - \boldsymbol{u} \rangle + \langle \boldsymbol{y}, \boldsymbol{u} \rangle - f(\boldsymbol{u}) - g(\boldsymbol{x} - \boldsymbol{u}) \right\} \\ &= \sup_{\boldsymbol{x}, \boldsymbol{u}} \left\{ \langle \boldsymbol{y}, \boldsymbol{x} - \boldsymbol{u} \rangle - g(\boldsymbol{x} - \boldsymbol{u}) \right\} + \sup_{\boldsymbol{u}} \left[\langle \boldsymbol{y}, \boldsymbol{u} \rangle - f(\boldsymbol{u}) \right] \\ &= g^*(\boldsymbol{y}) + f^*(\boldsymbol{y}) \end{aligned}$$

Recall \Box is similar to convolution: let \mathcal{F} denotes Fourier transform and let \star denotes integral convolution. Then $\mathcal{F}(f \star g) = \mathcal{F}(f) \cdot \mathcal{F}(g)$. Note the correspondence between $(+, \times)$ -algebra and $(\min, +)$ -algebra

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 $f\,\Box\,\frac{1}{2\mu}\|\cdot\|^2$

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Jean Moreau (1923-2014) Kosaku Yosida (1909-1990)

Moreau-Yosida envelope

- Moreau-Yosida envelope is the special case of infimal convolution $f \Box \frac{1}{2u} \| \cdot \|^2$
- Denote e_f^{μ} the Moreau-Yosida envelope of f under smoothing parameter $\mu > 0$,

$$e_f^{\mu}(\boldsymbol{x})\coloneqq \inf_{\boldsymbol{x}_1+\boldsymbol{x}_2=\boldsymbol{x}}f(\boldsymbol{x}_1)+rac{1}{2\mu}\|\boldsymbol{x}_2\|_2^2.$$

In optimization, usually the following form is used

$$e_f^{\mu}(m{x}) = \inf_{m{w}} f(m{w}) + rac{1}{2\mu} \|m{x} - m{w}\|_2^2 \stackrel{\|m{x} - m{w}\|_2^2 = \|m{w} - m{x}\|_2^2}{=} \inf_{m{w}} f(m{w}) + rac{1}{2\mu} \|m{w} - m{x}\|_2^2$$

The point-to-point map associated with Moreau-Yosida envelope is called the Proximal operator

$$\operatorname{prox}_f^{\mu}(\boldsymbol{x})\coloneqq \operatorname{argmin}_{\boldsymbol{w}} f(\boldsymbol{w}) + \frac{1}{2\mu}\|\boldsymbol{w} - \boldsymbol{x}\|_2^2$$

- Why study e_f^{μ} and $\operatorname{prox}_f^{\mu}$: they form the basis of modern convex optimization toolbox!
- Remark: inf becomes min if f is closed (epi f is closed) and convex

Why Moreau-Yosida envelope is useful

"Smoothing a non-smooth function to ease optimization"

- Consider minimizing f(x) = |x| using gradient descent $x^+ = x \alpha \nabla f(x)$
- Gradient descent requires differentiable f, while |x| is not differentiable at x = 0.
- Now instead of $\min f(x)$, consider $\min e_f^{\mu}(x)$, here is the magic:
 - e_f^{μ} is always differentiable

can use gradient descent!

• $\min e_f^{\mu}(x)$ and $\min f(x)$ share the same minimizer

Properties of Moreau-Yosida envelope and proximal operator

 $(\mu > 0)$

$$e_f^{\mu}$$
 is always differentiablesmoothness $\min e_f^{\mu}(x) = \min f(x)$ and $\operatorname{argmin} e_f^{\mu}(x) = \operatorname{argmin} f(x)$ same minimum and minimizer f is L -Lipschitz, then $0 \le f(x) - e_f^{\mu}(x) \le L^2 \mu$ e_f^{μ} is a lower bound of f $\nabla e_f^{\mu}(v) = \frac{1}{\mu} (v - \operatorname{prox}_f^{\mu}(v)) \in \partial f (\operatorname{prox}_f^{\mu}(v))$ relationship between e_f^{μ} and $\operatorname{prox}_f^{\mu}$ P roximal point algorithm = gradient descent on $\min e_f^{\mu}(x)$ Proximal point algorithm f is nonconvex, then $\operatorname{prox}_f^{\mu} : \mathbb{R}^n \Rightarrow \mathbb{R}^n$ non-uniqueness of $\operatorname{prox}_f^{\mu}$ f is convex, then $\operatorname{prox}_f^{\mu} : \mathbb{R}^n \to \mathbb{R}^n$ uniqueness of $\operatorname{prox}_f^{\mu}$ $e_f^{\mu} : \mathbb{R}^n \to \mathbb{R}$ uniqueness of e_f^{μ} $e_f^{\mu} : \mathbb{R}^n = \operatorname{argmin}_{g} f(y)$ is firmly non-expansive $\operatorname{Fix} \operatorname{prox}_f^{\mu} = \operatorname{argmin}_{g} f(\xi)$ weakly convergence $\operatorname{prox}_f^{\mu} (x) + \operatorname{prox}_{f^*}^{\mu} (x) = x$ Moreau decomposition

Example: Moreau-Yosida envelope of absolute value = Huber function

$$f(x) = \begin{cases} -x & x \le 0\\ x & x \ge 0 \end{cases}, \quad g = \frac{1}{2\mu} |\cdot|^2, \quad e_f^1(x) = (f \Box g)(x) = \begin{cases} -x - \frac{1}{2} & x < -1\\ \frac{1}{2}x^2 & x \in [-1, +1]\\ x - \frac{1}{2} & x > 1 \end{cases}$$

Example: $\min e_f^{\mu}(x) = \min f(x)$ and $\operatorname{argmin} e_f^{\mu}(x) = \operatorname{argmin} f(x)$

Proof by definition.

$$\begin{split} \min_{x} e_{f}^{\mu}(x) &= \min_{x} \left\{ \min_{y} \left\{ f(y) + \frac{1}{2\mu} \|x - y\|_{2}^{2} \right\} \right\} & \text{by definition of } e_{f}^{\mu} \\ &= \min_{y} \left\{ \min_{x} \left\{ f(y) + \frac{1}{2\mu} \|x - y\|_{2}^{2} \right\} \right\} & \text{you can swap the order of two min} \\ &= \min_{y} \left\{ \left\{ f(y) + \min_{x} \frac{1}{2\mu} \|x - y\|_{2}^{2} \right\} \right\} & f(y) \text{ is constant for } \min_{x} \\ &= \min_{y} f(y) & x = y \text{ minimizes } \|x - y\|_{2}^{2} \\ &= \min_{x} f(x) & \text{rename } y \text{ as } x \end{split}$$

Similar proof for argmin $e_f^{\mu}(x) = \operatorname{argmin} f(x)$.

Last page - summary

▶ The epigraph of
$$f : \mathbb{R}^n \to \overline{\mathbb{R}}$$
 is the set $\operatorname{epi} f \coloneqq \Big\{ (\boldsymbol{x}, \alpha) \in \mathbb{R}^n \times \mathbb{R} : \alpha \ge f(\boldsymbol{x}) \Big\}.$

Level sets

• Given $f_1 : \mathbb{R}^n \to \overline{\mathbb{R}}$ and $f_2 : \mathbb{R}^n \to \overline{\mathbb{R}}$. The infimal convolution $f_1 \Box f_2 : \mathbb{R}^n \to \overline{\mathbb{R}}$ is defined as

$$egin{array}{rll} (f_1 \Box f_2)(m{x}) &\coloneqq& \inf_{m{x}_1 + m{x}_2 = m{x}} f_1(m{x}_1) + f_2(m{x}_2) \ &=& \inf_{m{x}_1} & f_1(m{x}_1) + f_2(m{x} - m{x}_1) \ &=& \inf_{m{x}_2} & f_1(m{x} - m{x}_2) + f_2(m{x}_2) \end{array}$$

- Infimal convolution is also called epi-addition
- Moreau-Yosida envelope $f \Box \frac{1}{2\mu} \| \cdot \|^2$ as the foundation of modern optimization toolbox

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