

# Epigraphs, Infimal convolution and Moreau-Yosida envelope

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$$\text{epi} f := \left\{ (\mathbf{x}, \alpha) \in \mathbb{R}^n \times \mathbb{R} : \alpha \geq f(\mathbf{x}) \right\}.$$

# Graph and epigraph

► **(Graph of abstract function)**

Given two sets  $X$  (domain) and  $Y$  (codomain), the graph of the function  $f : X \rightarrow Y$  is the set

$$\text{graph } f := \{(\mathbf{x}, \mathbf{y}) \in X \times Y : \mathbf{y} = f(\mathbf{x})\}.$$

► **(Epigraph of function from  $\mathbb{R}^n$  to  $\overline{\mathbb{R}}$ )**

Let  $\overline{\mathbb{R}} := \mathbb{R} \cup \{+\infty\}$  be the extended reals.

Now consider  $X = \mathbb{R}^n$  and  $Y = \overline{\mathbb{R}}$ . For  $f : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$ , the epigraph of  $f$  is the set

$$\text{epi } f := \{(\mathbf{x}, \alpha) \in \mathbb{R}^n \times \mathbb{R} : \alpha \geq f(\mathbf{x})\}.$$

► **(Strict epigraph)**  $\text{epi}_S f = \text{epi } f \setminus \text{graph } f$ , or equivalently

$$\text{epi}_S f := \{(\mathbf{x}, \alpha) \in \mathbb{R}^n \times \mathbb{R} : \alpha > f(\mathbf{x})\}.$$

## Remarks

$$\text{epi}f := \left\{ (\mathbf{x}, \alpha) \in \mathbb{R}^n \times \mathbb{R} : \alpha \geq f(\mathbf{x}) \right\}$$

▶  $f : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$  but  $\text{epi}f$  is defined to be a subset of  $\mathbb{R}^n \times \mathbb{R}$ , not  $\mathbb{R}^n \times \overline{\mathbb{R}}$ .

▶ This is intentional to define  $\text{epi}f$  as a subset of  $\mathbb{R}^n \times \mathbb{R}$ .

▶  $\mathbb{R}^n$  is a vector space

▶  $\mathbb{R}^n \times \mathbb{R}$  is a vector space

▶  $\mathbb{R}^n \times \overline{\mathbb{R}}$  is **not** a vector space:  $\nexists$  additive identity for  $\infty + \infty$

(Being a vector space allows to use tools from real analysis and functional analysis.)

▶ **(At infinity)** If  $f(\mathbf{x}_0) = +\infty$  at  $\mathbf{x} = \mathbf{x}_0$ , then  $(\mathbf{x}_0, +\infty) \notin \text{epi}f$ .

▶ **Two extreme cases**

▶ If  $f_{\infty+}(\mathbf{x}) = +\infty$  for all  $\mathbf{x}$ , then  $\text{epi}f_{\infty+}$  is  $\emptyset$ .

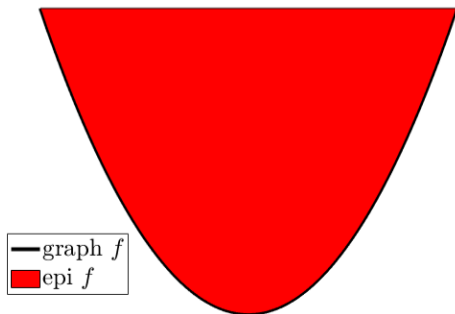
▶ If  $f_{\infty-}(\mathbf{x}) = -\infty$  for all  $\mathbf{x}$ , then  $\text{epi}f_{\infty-}$  is the whole  $\mathbb{R}^n \times \mathbb{R}$

(Convention:  $-\infty < +\infty$ )

**Empty epigraph**  
**Whole space epigraph**

## Visualization of $\text{graph } f$ and $\text{epi } f$

- ▶  $\text{epi } f$  consists of **all** the points of  $\mathbb{R}^{n+1}$  lying **on or above**  $\text{graph } f$ .
- ▶ Example:  $f(x) = x^2$ 
  - ▶  $n = 1$  (1-dimensional)
  - ▶  $\text{graph } f := \{(x, y) \in \mathbb{R} \times \mathbb{R} : y = f(x)\}$  is a 1-dimensional curve in a 2-dimensional space.
  - ▶  $\text{epi } f := \{(x, \alpha) \in \mathbb{R} \times \mathbb{R} : \alpha \geq f(x)\}$  is a 2-dimensional set in a 2-dimensional space.



## Level sets: a concept related to epigraph.

$$\text{lev}_{\leq \alpha} f := \{ \mathbf{x} \in \mathbb{R}^n : f(\mathbf{x}) \leq \alpha \}$$

$$\text{lev}_{< \alpha} f := \{ \mathbf{x} \in \mathbb{R}^n : f(\mathbf{x}) < \alpha \}$$

► For  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  and  $\alpha \in \mathbb{R}$ , define  $\text{lev}_{= \alpha} f := \{ \mathbf{x} \in \mathbb{R}^n : f(\mathbf{x}) = \alpha \}$

$$\text{lev}_{> \alpha} f := \{ \mathbf{x} \in \mathbb{R}^n : f(\mathbf{x}) > \alpha \}$$

$$\text{lev}_{\geq \alpha} f := \{ \mathbf{x} \in \mathbb{R}^n : f(\mathbf{x}) \geq \alpha \}$$

The important one for minimization is  $\text{lev}_{\leq \alpha} f$  (named sublevel sets).

►  $\text{lev}$  is a subset of domain, not codomain.

► If  $\alpha = \inf f$ , then  $\text{lev}_{\leq \alpha} f = \text{lev}_{= \alpha} f = \text{argmin } f$ .

► Level sets can be

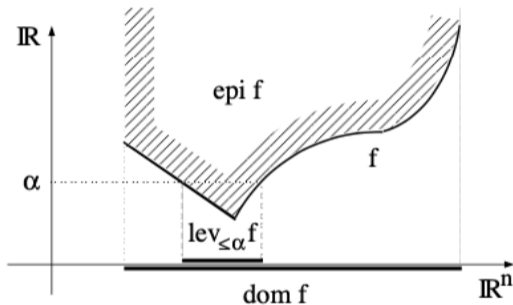
► empty:  $\text{lev}_{\leq -1}(x^2)$ , no  $x$  makes  $x^2 \leq -1$

► non-continuous:  $\text{lev}_{=0} \sin(x)$ , then  $x = \{n\pi\}_{n \in \mathbb{N}}$  is a set of dots

► non-convex:  $\text{lev}_{=0} \sin(x)$

# An illustration of $\text{lev}_{\leq \alpha} f$

From the book by Rockafellar and Wets.



Rockafellar, R. Tyrrell, and Roger J-B. Wets. Variational analysis. Springer, 2009.

## Q & A

Optimization newbie: “Why do we talk about epigraph?”

Optimization expert: “Because it is useful!”

The main philosophy

- ▶ Many properties of  $f$  has a counterpart in  $\text{epi}f$ .
- ▶ Sometimes it is easier to work on  $\text{epi}f$  than on  $f$ .



## Properties of $f$

- ▶  $f$  is proper
- ▶  $f$  is closed
- ▶  $f$  is lower semicontinuous on  $\mathbb{R}^n$
- ▶  $f$  is convex
- ▶  $f$  is strictly convex
- ▶  $f \square g$  (Infimal convolution of  $f$  and  $g$ )

## Properties of $\text{epi} f$

- $\text{epi} f$  is nonempty
- $\text{epi} f$  is a closed set
- $\text{epi} f$  is nonempty and closed in  $\mathbb{R}^n \times \mathbb{R}$
- $\text{epi} f$  is a convex set
- $\text{epi}_S f$  is a convex set
- Minkowski sum of  $\text{epi} f$  and  $\text{epi} g$

## Example: $f$ is convex $\iff$ $\text{epi}f$ is convex

(To avoid confusing the bracket of graph and bracket of a number, here we use  $\{\mathbf{x}_1, y_1\} \in \text{epi}f$ )

► ( $\implies$ )

►  $f$  is convex  $\iff f((1-t)\mathbf{x}_1 + t\mathbf{x}_2) \leq (1-t)f(\mathbf{x}_1) + tf(\mathbf{x}_2) \forall \mathbf{x}_1, \mathbf{x}_2$  and  $t \in [0, 1]$ . (\*)

►  $\{\mathbf{x}_1, y_1\} \in \text{epi}f \iff f(\mathbf{x}_1) \leq y_1$  and  $\{\mathbf{x}_2, y_2\} \in \text{epi}f \iff f(\mathbf{x}_2) \leq y_2$  (\*\*)

► Now  $f((1-t)\mathbf{x}_1 + t\mathbf{x}_2) \stackrel{(*)}{\leq} (1-t)f(\mathbf{x}_1) + tf(\mathbf{x}_2) \stackrel{(**)}{\leq} (1-t)y_1 + ty_2$ . So

$\{(1-t)\mathbf{x}_1 + t\mathbf{x}_2, (1-t)y_1 + ty_2\} = (1-t)\{\mathbf{x}_1, y_1\} + t\{\mathbf{x}_2, y_2\} \stackrel{(**)}{\in} \text{epi}f$ . So  $\text{epi}f$  is a convex set.

► ( $\impliedby$ )

►  $\text{epi}f$  is a convex set  $\iff \{\mathbf{x}_1, y_1\} \in \text{epi}f, \{\mathbf{x}_2, y_2\} \in \text{epi}f$  implies for  $t \in [0, 1]$  we have

$$(1-t)\{\mathbf{x}_1, y_1\} + t\{\mathbf{x}_2, y_2\} = \{(1-t)\mathbf{x}_1 + t\mathbf{x}_2, (1-t)y_1 + ty_2\} \in \text{epi}f. \quad (***)$$

► By definition of epigraph, (\*\*\*) is equivalent to

$$f(\mathbf{x}_1) \leq y_1 \text{ and } f(\mathbf{x}_2) \leq y_2 \text{ implies } f((1-t)\mathbf{x}_1 + t\mathbf{x}_2) \leq (1-t)y_1 + ty_2.$$

Choose  $f(\mathbf{x}_1) \leq y_1$  and  $f(\mathbf{x}_2) \leq y_2$  gives  $f((1-t)\mathbf{x}_1 + t\mathbf{x}_2) \leq (1-t)f(\mathbf{x}_1) + tf(\mathbf{x}_2)$ , so  $f$  is convex.

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$$(f_1 \square f_2)(\mathbf{x}) := \inf_{\mathbf{x}_1 + \mathbf{x}_2 = \mathbf{x}} f_1(\mathbf{x}_1) + f_2(\mathbf{x}_2)$$

# Infimal convolution

- ▶ Given  $f_1 : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$  and  $f_2 : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$ . The infimal convolution  $f_1 \square f_2 : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$  is defined as

$$\begin{aligned}(f_1 \square f_2)(\mathbf{x}) &:= \inf_{\mathbf{x}_1 + \mathbf{x}_2 = \mathbf{x}} f_1(\mathbf{x}_1) + f_2(\mathbf{x}_2) \\ &= \inf_{\mathbf{x}_1} f_1(\mathbf{x}_1) + f_2(\mathbf{x} - \mathbf{x}_1) \\ &= \inf_{\mathbf{x}_2} f_1(\mathbf{x} - \mathbf{x}_2) + f_2(\mathbf{x}_2)\end{aligned}$$

Convention:  $\infty - \infty = \infty$  and  $\inf \emptyset = +\infty$

- ▶ History

- ▶ Earliest(?) work  
Hausdorff, Über halbstetige Funktionen und deren Verallgemeinerung, Math. Zeit. 5 (1919), 292-309.
- ▶ Fenchel, "Convex Cones, Sets, and Functions", Lecture Notes, Princeton University, Princeton, 1953.
- ▶ First systematic study of infimal convolution  
Moreau, Inf-convolution, Sémin. d'Math. Montpellier (1963), 3.1-3.48
- ▶ Later works by Attouch, Rockafellar, Hiriart-Urruty, etc
- ▶ Thomas Stromberg's PhD thesis (1994): a nice summary.

# How infimal convolution gets its name

$$(f_1 \square f_2)(\mathbf{x}) := \inf_{\mathbf{x}_1} f_1(\mathbf{x}_1) + f_2(\mathbf{x} - \mathbf{x}_1).$$

- ▶ Definition of (integral) convolution (Examples: Laplace transform, Fourier transform.)

$$(f * g)(t) := \int_{-\infty}^{\infty} f(\tau)g(t - \tau)d\tau.$$

- ▶  $f_1(\mathbf{x}_1) + f_2(\mathbf{x} - \mathbf{x}_1)$  “looks similar” to (integral) convolution  $\implies$  people name it convolution.

- ▶ There is  $\inf_w \implies$  people name it infimal.

- ▶ Deep fact: integral convolution is in  $(+, \times)$ -algebra

- ▶ Integration = summation
- ▶ You combine  $f$  and  $g$  by multiplication

Infimal convolution is in  $(\min, +)$ -algebra (it forms a tropical semi-ring)

- ▶ The summation is replaced by  $\min$
- ▶ You combine  $f$  and  $g$  by addition

# What infimal convolution solves: an economics example

- ▶ You want to buy  $n$  hamburgers from MacDonal'd and Burger King.  
Suppose buying  $n_1$  hamburgers from MacDonal'd costs you  $f(n_1)$ , and if you buy  $n_2$  hamburgers from Burger King, the price is  $g(n_2)$ .
- ▶ You want to find the infimum of the total cost  $f(n_1) + g(n_2)$  subject to the constraint  $n_1 + n_2 = n$ . I.e., you want to find the “cheapest way” to buy  $n$  hamburgers.
- ▶ This problem is exactly: calculate  $(f \square g)(n)$

$$(f \square g)(n) = \inf_{n_1+n_2=n} f(n_1) + g(n_2) = \underbrace{\inf_{n_1} f(n_1) + g(n - n_1)}_{\text{focus on } n_1} = \underbrace{\inf_{n_2} f(n - n_2) + g(n_2)}_{\text{focus on } n_2}.$$

this also means infimal convolution is commutative

Infimal convolution is commutative:  $f \square g = g \square f$ .

## Example. Infimal convolution of two indicator functions

$$\begin{aligned}(f_1 \square f_2)(\mathbf{x}) &:= \inf_{\mathbf{x}_1 + \mathbf{x}_2 = \mathbf{x}} f_1(\mathbf{x}_1) + f_2(\mathbf{x}_2) \\ &= \inf_{\mathbf{x}_1} f_1(\mathbf{x}_1) + f_2(\mathbf{x} - \mathbf{x}_1) \\ &= \inf_{\mathbf{x}_2} f_1(\mathbf{x} - \mathbf{x}_2) + f_2(\mathbf{x}_2)\end{aligned}$$

- ▶ Given two sets  $C_1, C_2$  and two indicator functions  $i_{C_1}, i_{C_2}$ .

$$(i_{C_1} \square i_{C_2})(\mathbf{x}) = \inf_{\mathbf{x}_1 + \mathbf{x}_2 = \mathbf{x}} i_{C_1}(\mathbf{x}_1) + i_{C_2}(\mathbf{x}_2) = i_{C_1 \oplus C_2}(\mathbf{x})$$

$\oplus$  is Minkowski sum of sets:  $P \oplus Q := \{p + q \mid p \in P, q \in Q\}$ .

- ▶ As convexity of set conserves under Minkowski sum, so  $C_1 \oplus C_2$  is a convex set and  $i_{C_1 \oplus C_2}$  is a convex function. Here we see that inf-convolution of two convex functions is a convex function.
- ▶ In general, if  $f_1 : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$  and  $f_2 : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$  are convex functions, then  $f_1 \square f_2$  is also convex.  
Proof. We prove inf-convolution preserves convexity using definition and operations that preserve convexity. By definition,  $(f_1 \square f_2)(x) = \inf_{x_1} h(x, x_1)$  where  $h(x, x_1) = f_1(x_1) + f_2(x - x_1)$ . By assumption  $f_1(x_1)$  is convex and  $f_2(x_1)$  is convex. The function  $f_2(x - x_1)$  is  $f_2(x_1)$  with argument  $x_1$  under a translation to  $x - x_1$  so  $f_2(x - x_1)$  is convex. Now  $h(x, x_1)$  is the sum of two convex functions on  $x_1$ , thus it is convex.

## Infimal convolution is also called epi-addition

$$(f_1 \square f_2)(\mathbf{x}) := \inf_{\mathbf{x}_2} f_1(\mathbf{x} - \mathbf{x}_2) + f_2(\mathbf{x}_2).$$

- ▶  $\text{epi}_S(f_1 \square f_2) = \text{epi}_S f_1 \oplus \text{epi}_S f_2 \iff$  inf-convolution of convex functions is convex<sup>1</sup>
- ▶  $\text{epi}(f_1 \square f_2) \supseteq \text{epi} f_1 \oplus \text{epi} f_2$
- ▶  $\text{epi}(f_1 \square f_2) = \text{epi} f_1 \oplus \text{epi} f_2$  if inf-convolution is exact  
Exact means the inf is gone:  $(f_1 \square f_2)(\mathbf{x}) = f_1(\mathbf{x} - \mathbf{x}_2) + f_2(\mathbf{x}_2)$ .
- ▶ For proof, see Jean Jacques Moreau. Inf-convolution, sous-additivité, convexité des fonctions numériques. Journal de Mathématiques Pures et Appliquées, 1970.  
<https://hal.archives-ouvertes.fr/hal-02162006>

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<sup>1</sup>Remark 2.3.3 in Urruty, Jean-Baptiste Hiriart, and Claude Lemaréchal. Convex analysis and minimization algorithms. Springer-Verlag, 1993



## The proof of $\text{epi}(f_1 \square f_2) \supseteq \text{epi}f_1 \oplus \text{epi}f_2$

- ▶ Take  $\{\mathbf{x}, \alpha\} \in \text{epi}f_1 \oplus \text{epi}f_2$ . Since the element  $\{\mathbf{x}, \alpha\}$  is inside the sum of two sets ( $\text{epi}f_1$  and  $\text{epi}f_2$ ), that means we can decompose  $\{\mathbf{x}, \alpha\}$  as the sum of element from each set. I.e., we have  $\mathbf{x} = \mathbf{x}_1 + \mathbf{x}_2$  and  $\alpha = \alpha_1 + \alpha_2$  that  $\{\mathbf{x}_1, \alpha_1\} \in \text{epi}f_1$  and  $\{\mathbf{x}_2, \alpha_2\} \in \text{epi}f_2$ .  
(It means given  $\{\mathbf{x}, \alpha\}$ , there is exist the pair  $\{\mathbf{x}_1, \alpha_1\}$ ,  $\{\mathbf{x}_2, \alpha_2\}$  that fulfill the above conditions)

- ▶  $\alpha_1$  implies  $f_1(\mathbf{x}_1) \leq \alpha_1$ ,  $\alpha_2$  implies  $f_2(\mathbf{x}_2) \leq \alpha_2$  and  $f_1(\mathbf{x}_1) + f_2(\mathbf{x}_2) \leq \alpha_1 + \alpha_2 = \alpha$ .

- ▶ Now consider  $(f_1 \square f_2)(\mathbf{x}) := \inf_{\mathbf{y}_1 + \mathbf{y}_2 = \mathbf{x}} f_1(\mathbf{y}_1) + f_2(\mathbf{y}_2)$ . As  $\mathbf{x}_1 + \mathbf{x}_2 = \mathbf{x} = \mathbf{y}_1 + \mathbf{y}_2$ , the infimum  $\inf_{\mathbf{y}_1 + \mathbf{y}_2 = \mathbf{x}}$  is the smallest among all pair that sum to  $\mathbf{x}$ , so

$$(f_1 \square f_2)(\mathbf{x}) := \inf_{\mathbf{y}_1 + \mathbf{y}_2 = \mathbf{x}} f_1(\mathbf{y}_1) + f_2(\mathbf{y}_2) \leq f_1(\mathbf{x}_1) + f_2(\mathbf{x}_2) \leq \alpha.$$

So  $\{\mathbf{x}, \alpha\} \in \text{epi}(f_1 \square f_2)$ .

- ▶ What we just showed is  $\{\mathbf{x}, \alpha\} \in \text{epi}f_1 \oplus \text{epi}f_2 \implies \{\mathbf{x}, \alpha\} \in \text{epi}(f_1 \square f_2)$ , so in set language we have  $\text{epi}(f_1 \square f_2) \supseteq \text{epi}f_1 \oplus \text{epi}f_2$ .

## Example. Pictorial illustration of epi-addition (and Minkowski sum)

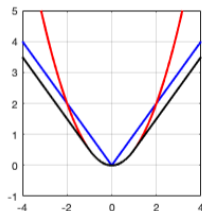
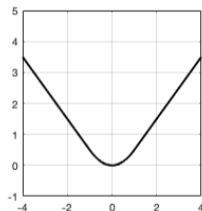
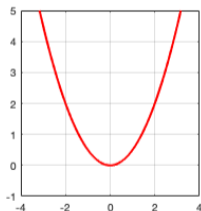
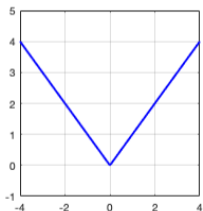
►  $f = |\cdot|$

$$\text{epi}f = \{(x, t) \in \mathbb{R} \times \mathbb{R} : |x| \leq t\}$$

►  $g = \frac{1}{2}(\cdot)^2$

$$\text{epi}g = \{(x, t) \in \mathbb{R} \times \mathbb{R} : \frac{1}{2}x^2 \leq t\}$$

►  $(f \square g)(x) = \inf_w |w| + \frac{1}{2}(x - w)^2$        $\text{epi}(f \square g) = \{(x, t) \in \mathbb{R} \times \mathbb{R} : \left(\inf_w |w| + \frac{1}{2}(x - w)^2\right) \leq t\}$



“Epi-addition: sliding the blue curve on red curve and perform union operation gives the black curve”

## What about $f_1 \square f_2 \square f_3$ ? Inf-convolution is associative

$$\begin{aligned} f_1 \square (f_2 \square f_3)(t) &= \inf_{x+y=t} \left\{ f_1(x) + (f_2 \square f_3)(y) \right\} && \text{by definition} \\ &= \inf_{x+y=t} \left\{ f_1(x) + \left\{ \inf_{z+w=y} f_2(w) + f_3(z) \right\} \right\} && \text{by definition} \\ &= \inf_{\substack{x+y=t \\ z+w=y}} \left\{ f_1(x) + f_2(w) + f_3(z) \right\} && \text{you can move inf} \\ &= \inf_{x+z+w=t} \left\{ f_1(x) + f_2(w) + f_3(z) \right\} && \text{combine } x + y = t, z + w = y \\ &= \inf_{\substack{r+z=t \\ x+w=r}} \left\{ f_1(x) + f_2(w) + f_3(z) \right\} && \text{let } t = r + z, x + w = r \\ &= \inf_{r+z=t} \left\{ \left\{ \inf_{x+w=r} f_1(x) + f_2(w) \right\} + f_3(z) \right\} && \text{you can move inf} \\ &= \inf_{r+z=t} \left\{ (f_1 \square f_2)(r) + f_3(z) \right\} && \text{by definition} \\ &= (f_1 \square f_2) \square f_3(t) && \text{by definition} \end{aligned}$$

# Properties of inf-convolution

We already see

- ▶  $f \square g = g \square f$  commutative
- ▶  $f \square g \square h = (f \square g) \square h = f \square (g \square h)$  associative
- ▶  $f, g$  convex  $\implies f \square g$  convex inf-convolution preserves convexity

Useful table

$f$	$g$	$f \square g$
$f$	0	$\inf_x f(x)$
$i_C$	$\ \cdot\ _2$	$d_C$
$i_{C_1}$	$i_{C_2}$	$i_{C_1 \oplus C_2}$
$f$	$i_x$	$f(\cdot - x)$
$f$	$\langle s, \cdot \rangle$	$\langle s, \cdot \rangle - f^*(s)$
$f$ convex	$f$	$2f(\frac{\cdot}{2})$

- ▶ Distance function  $d_C(x) = \inf_{c \in C} \|x - c\|_2$
- ▶ Indicator function  $i_C(x) = \begin{cases} 0 & x \in C \\ +\infty & x \notin C \end{cases}$
- ▶ Conjugate  $f^*(x) = \sup_u \langle u, x \rangle - f(x)$

# Inf-convolution and conjugate: $(f \square g)^* = f^* + g^*$

Let  $f, g : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$  be proper functions.

$$\begin{aligned}(f \square g)^*(\mathbf{y}) &= \sup_{\mathbf{x}} \left\{ \langle \mathbf{y}, \mathbf{x} \rangle - (f \square g)(\mathbf{x}) \right\} && \text{by definition of conjugate} \\ &= \sup_{\mathbf{x}} \left\{ \langle \mathbf{y}, \mathbf{x} \rangle - \inf_{\mathbf{u}} [f(\mathbf{u}) + g(\mathbf{x} - \mathbf{u})] \right\} && \text{by definition of inf-convolution} \\ &= \sup_{\mathbf{x}} \sup_{\mathbf{u}} \left\{ \langle \mathbf{y}, \mathbf{x} \rangle - [f(\mathbf{u}) + g(\mathbf{x} - \mathbf{u})] \right\} && - \inf_{\mathbf{u}} = + \sup_{\mathbf{u}} \\ &= \sup_{\mathbf{x}, \mathbf{u}} \left\{ \langle \mathbf{y}, \mathbf{x} - \mathbf{u} + \mathbf{u} \rangle - f(\mathbf{u}) - g(\mathbf{x} - \mathbf{u}) \right\} \\ &= \sup_{\mathbf{x}, \mathbf{u}} \left\{ \langle \mathbf{y}, \mathbf{x} - \mathbf{u} \rangle + \langle \mathbf{y}, \mathbf{u} \rangle - f(\mathbf{u}) - g(\mathbf{x} - \mathbf{u}) \right\} \\ &= \sup_{\mathbf{x}, \mathbf{u}} \left\{ \langle \mathbf{y}, \mathbf{x} - \mathbf{u} \rangle - g(\mathbf{x} - \mathbf{u}) \right\} + \sup_{\mathbf{u}} [\langle \mathbf{y}, \mathbf{u} \rangle - f(\mathbf{u})] \\ &= g^*(\mathbf{y}) + f^*(\mathbf{y})\end{aligned}$$

Recall  $\square$  is similar to convolution: let  $\mathcal{F}$  denotes Fourier transform and let  $\star$  denotes integral convolution. Then  $\mathcal{F}(f \star g) = \mathcal{F}(f) \cdot \mathcal{F}(g)$ . Note the correspondence between  $(+, \times)$ -algebra and  $(\min, +)$ -algebra

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$$f \square \frac{1}{2\mu} \|\cdot\|^2$$

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Jean Jacques Moreau (1923-2014)



Kosaku Yosida (1909-1990)

## Moreau-Yosida envelope

- ▶ Moreau-Yosida envelope is the special case of infimal convolution  $f \square_{\frac{1}{2\mu}} \|\cdot\|^2$
- ▶ Denote  $e_f^\mu$  the Moreau-Yosida envelope of  $f$  under smoothing parameter  $\mu > 0$ ,

$$e_f^\mu(\mathbf{x}) := \inf_{\mathbf{x}_1 + \mathbf{x}_2 = \mathbf{x}} f(\mathbf{x}_1) + \frac{1}{2\mu} \|\mathbf{x}_2\|_2^2.$$

In optimization, usually the following form is used

$$e_f^\mu(\mathbf{x}) = \inf_{\mathbf{w}} f(\mathbf{w}) + \frac{1}{2\mu} \|\mathbf{x} - \mathbf{w}\|_2^2 \quad \|\mathbf{x} - \mathbf{w}\|_2^2 = \|\mathbf{w} - \mathbf{x}\|_2^2 \quad \inf_{\mathbf{w}} f(\mathbf{w}) + \frac{1}{2\mu} \|\mathbf{w} - \mathbf{x}\|_2^2$$

- ▶ The point-to-point map associated with Moreau-Yosida envelope is called the Proximal operator

$$\text{prox}_f^\mu(\mathbf{x}) := \underset{\mathbf{w}}{\text{argmin}} f(\mathbf{w}) + \frac{1}{2\mu} \|\mathbf{w} - \mathbf{x}\|_2^2$$

- ▶ Why study  $e_f^\mu$  and  $\text{prox}_f^\mu$ : they form the basis of modern convex optimization toolbox!
- ▶ Remark:  $\inf$  becomes  $\min$  if  $f$  is closed ( $\text{epi} f$  is closed) and convex

# Why Moreau-Yosida envelope is useful

“Smoothing a non-smooth function to ease optimization”

- ▶ Consider minimizing  $f(x) = |x|$ .
- ▶ For general problem  $\min f(x)$ , a popular approach is gradient descent  $x^+ = x - \alpha \nabla f(x)$
- ▶ Using gradient descent requires  $f$  is differentiable, where  $f(x) = |x|$  is not differentiable at  $x = 0$ .
- ▶ Now instead of  $\min f(x)$ , consider  $\min e_f^\mu(x)$ , here is the magic:
  - ▶  $e_f^\mu$  is always differentiable can use gradient descent!
  - ▶  $\min e_f^\mu(x)$  and  $\min f(x)$  share the same minimizer



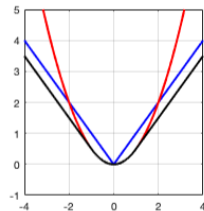
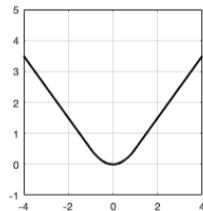
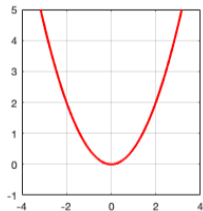
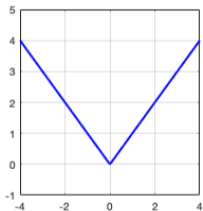
# Properties of Moreau-Yosida envelope and proximal operator

( $\mu > 0$ )

- ▶  $e_f^\mu$  is always differentiable smoothness
- ▶  $\min e_f^\mu(x) = \min f(x)$  and  $\operatorname{argmin} e_f^\mu(x) = \operatorname{argmin} f(x)$  same minimum and minimizer
- ▶  $f$  is  $L$ -Lipschitz, then  $0 \leq f(x) - e_f^\mu(x) \leq L^2\mu$   $e_f^\mu$  is a lower bound of  $f$
- ▶  $\nabla e_f^\mu(v) = \frac{1}{\mu}(v - \operatorname{prox}_f^\mu(v)) \in \partial f(\operatorname{prox}_f^\mu(v))$  relationship between  $e_f^\mu$  and  $\operatorname{prox}_f^\mu$
- ▶ Proximal point algorithm = gradient descent on  $\min e_f^\mu(x)$  Proximal point algorithm
- ▶  $f$  is nonconvex, then  $\operatorname{prox}_f^\mu : \mathbb{R}^n \rightrightarrows \mathbb{R}^n$  non-uniqueness of  $\operatorname{prox}_f^\mu$
- ▶  $f$  is convex, then  $\operatorname{prox}_f^\mu : \mathbb{R}^n \rightarrow \mathbb{R}^n$  uniqueness of  $\operatorname{prox}_f^\mu$
- ▶  $e_f^\mu : \mathbb{R}^n \rightarrow \mathbb{R}$  uniqueness of  $e_f^\mu$
- ▶  $\langle \operatorname{prox}_f^\mu(x) - \operatorname{prox}_f^\mu(y), (\operatorname{Id} - \operatorname{prox}_f^\mu)(x) - (\operatorname{Id} - \operatorname{prox}_f^\mu)(y) \rangle \geq 0$   $\operatorname{prox}_f^\mu$  is firmly non-expansive
- ▶ Fix  $\operatorname{prox}_f^\mu = \operatorname{argmin}_y f(y)$  fixed point
- ▶ Let  $T = \operatorname{prox}_f^\mu$ , then  $\{T^k x\}_{k \in \mathbb{N}} \rightarrow \operatorname{argmin}_\xi f(\xi)$  weakly convergence
- ▶  $\operatorname{prox}_f^\mu(x) + \operatorname{prox}_{f^*}^\mu(x) = x$  Moreau decomposition

# Example: Moreau-Yosida envelope of absolute value = Huber function

$$f(x) = \begin{cases} -x & x \leq 0 \\ x & x \geq 0 \end{cases}, \quad g = \frac{1}{2\mu} | \cdot |^2, \quad e_f^1(x) = (f \square g)(x) = \begin{cases} -x - \frac{1}{2} & x < -1 \\ \frac{1}{2}x^2 & x \in [-1, +1] \\ x - \frac{1}{2} & x > 1 \end{cases}$$



Example:  $\min e_f^\mu(x) = \min f(x)$  and  $\operatorname{argmin} e_f^\mu(x) = \operatorname{argmin} f(x)$

Proof by definition.

$$\begin{aligned}\min_x e_f^\mu(x) &= \min_x \left\{ \min_y \left\{ f(y) + \frac{1}{2\mu} \|x - y\|_2^2 \right\} \right\} && \text{by definition of } e_f^\mu \\ &= \min_y \left\{ \min_x \left\{ f(y) + \frac{1}{2\mu} \|x - y\|_2^2 \right\} \right\} && \text{you can swap the order of two min} \\ &= \min_y \left\{ \left\{ f(y) + \min_x \frac{1}{2\mu} \|x - y\|_2^2 \right\} \right\} && f(y) \text{ is constant for } \min_x \\ &= \min_y \left\{ f(y) \right\} && x = y \text{ minimizes } \|x - y\|_2^2 \\ &= \min_x f(x) && \text{rename } y \text{ as } x\end{aligned}$$

Similar proof for  $\operatorname{argmin} e_f^\mu(x) = \operatorname{argmin} f(x)$ .

## Last page - summary

- ▶ The epigraph of  $f : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$  is the set  $\text{epi} f := \left\{ (\mathbf{x}, \alpha) \in \mathbb{R}^n \times \mathbb{R} : \alpha \geq f(\mathbf{x}) \right\}$ .
- ▶ Level sets
- ▶ Given  $f_1 : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$  and  $f_2 : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$ . The infimal convolution  $f_1 \square f_2 : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$  is defined as

$$\begin{aligned}(f_1 \square f_2)(\mathbf{x}) &:= \inf_{\mathbf{x}_1 + \mathbf{x}_2 = \mathbf{x}} f_1(\mathbf{x}_1) + f_2(\mathbf{x}_2) \\ &= \inf_{\mathbf{x}_1} f_1(\mathbf{x}_1) + f_2(\mathbf{x} - \mathbf{x}_1) \\ &= \inf_{\mathbf{x}_2} f_1(\mathbf{x} - \mathbf{x}_2) + f_2(\mathbf{x}_2)\end{aligned}$$

- ▶ Infimal convolution is also called epi-addition
- ▶ Moreau-Yosida envelope  $f \square \frac{1}{2\mu} \|\cdot\|^2$  as the foundation of modern optimization toolbox

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