

# Shannon-Fano Code

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## Example

Given a set of symbols and their probabilities of occurrence

$$S = \{a, b, c, d, e\} \quad P(S) = \{0.4, 0.3, 0.15, 0.1, 0.05\}$$

First find the entropy (average information content per message)

$$H = \sum p_i I_i = \sum p_i \log_2 \frac{1}{p_i} = 2.0087 \text{ bit/symbol}$$

If normal binary code is used

<i>a</i>	000
<i>b</i>	001
<i>c</i>	010
<i>d</i>	011
<i>e</i>	100

Average code length      3

Therefore, we use a 3 bit code word to transmit 2.0087 bit of information (on average), the efficiency and redundancy are

$$\eta = \frac{H}{L_{avg}} = \frac{2.0087}{3} = 67\% \quad \gamma = 1 - \eta = 33\%$$

Now consider **Shannon-Fano Code**

The idea of Shannon-Fano Code is to first group the symbol into 2 group with equal probabilities ( or as close as possible )

$$a_{0.4}, b_{0.3}, c_{0.15}, d_{0.1}, e_{0.05}$$

The grouping is

$$\underbrace{a_{0.4}, d_{0.1}}_{0.5} \quad \underbrace{b_{0.3}, c_{0.15}, e_{0.05}}_{0.5}$$

Then assign the first bit

$$a_{0.4}(0), d_{0.1}(0) \quad b_{0.3}(1), c_{0.15}(1), e_{0.05}(1)$$

Repeat the step, the new grouping is

$$a_{0.4}(0) \quad d_{0.1}(0) \quad b_{0.3}(1) \quad c_{0.15}(1), e_{0.05}(1)$$

Assign next bit ( notice that although *b* is separated , it still need to assign one more bit because of pre-fix code condition ! )

$$a_{0.4}(00) \quad d_{0.1}(01) \quad b_{0.3}(10) \quad c_{0.15}(11), e_{0.05}(11)$$

Lastly split  $c$  and  $e$

$$a_{0.4}(00) \quad d_{0.1}(01) \quad b_{0.3}(10) \quad c_{0.15}(110) \quad e_{0.05}(111)$$

Done

The average code length is thus

$$L_{av} = \sum p_i l_i = 0.4(2) + 0.1(2) + 0.3(2) + 0.15(3) + 0.05(3) = 2.2 > H = 2.0087$$

The efficiency now is

$$\eta = \frac{H}{L_{av}} = \frac{2.0087}{2.2} = 91\% \quad \gamma = 1 - \eta = 9\%$$

Efficiency increased from 66% to 91%, improved a lot !

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